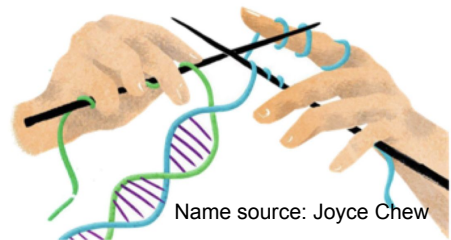


# Cal-DISKS

Calvin/Berkeley Distributed  $k$ -mer Sketcher

Elizabeth Koning

Advised by Joel Adams

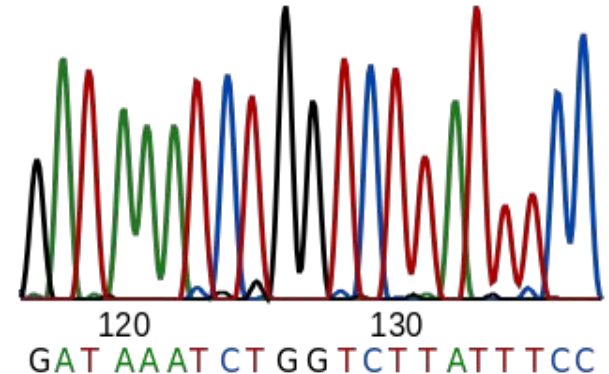


Name source: Joyce Chew

# Genetic Sequences

... are strings of A, C, G, T characters representing DNA

- Used in Computational Biology problems, including:
  - Metagenomics (DNA of many different organisms)
  - Genome assembly (properly ordering genomic sequences)
- Involves massive amounts of data
  - Human genome: 3 billion base pairs (>750 megabytes)
  - Often many terabytes of data for metagenomes



# *k*-mer Counting

**AGTTGAGTTGAG**

AGTTGA  
GTTGAG  
TTGAGT  
TGAGTT  
GAGTTG  
AGTTGA  
GTTGAG



{  
AGTTGA  
AGTTGA  
GTTGAG  
GTTGAG  
TTGAGT  
TGAGTT  
GAGTTG  
}

# k-mer Usage

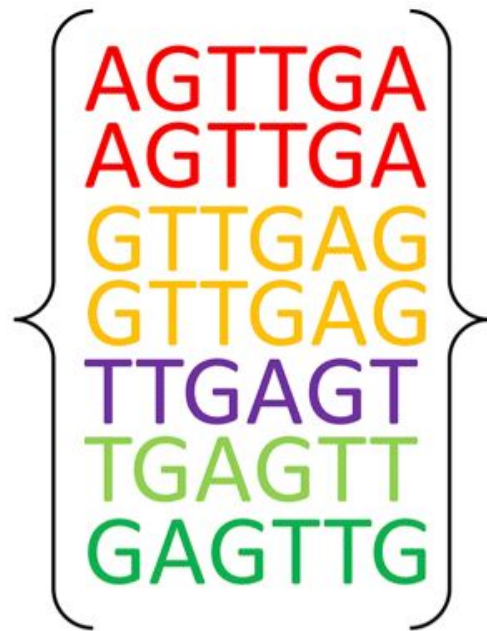
How similar are two DNA sequences?

Two applications:

1. *Alignment*: Find initial k-mer matches before performing nucleotide-level comparison of the genetic sequences. (If you need very accurate analysis)
2. *Similarity*: Computing the percentage of similarity between two sequences without performing computationally-intensive nucleotide-level alignment. (If you just need a rough estimation)

Problem for both applications:

We have many, many k-mers



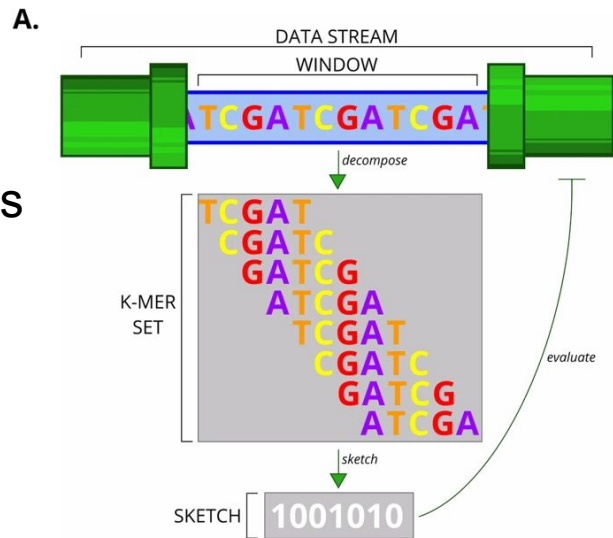
# k-mer Sketcher

Strategy to reduce the number of k-mers:

- Select a subset of a sequence's k-mers  
→ The subset is a “sketch” of its sequence and is orders of magnitude smaller than the full representation

Tradeoff:

- Smaller sketch → less storage and processing time
- Larger sketch → higher accuracy



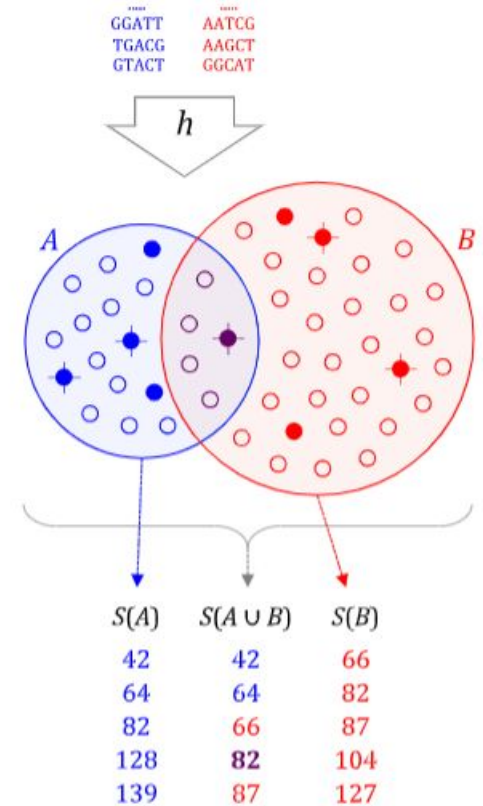
# Sketching Technique

Use MinHash data structure

- Use a hash function to hash each k-mer
- Store set of minimum k-mer hash-values

Benefit:

- Sketches are smaller to store and faster to compare to other sketches
- Results are more accurate than a random sample



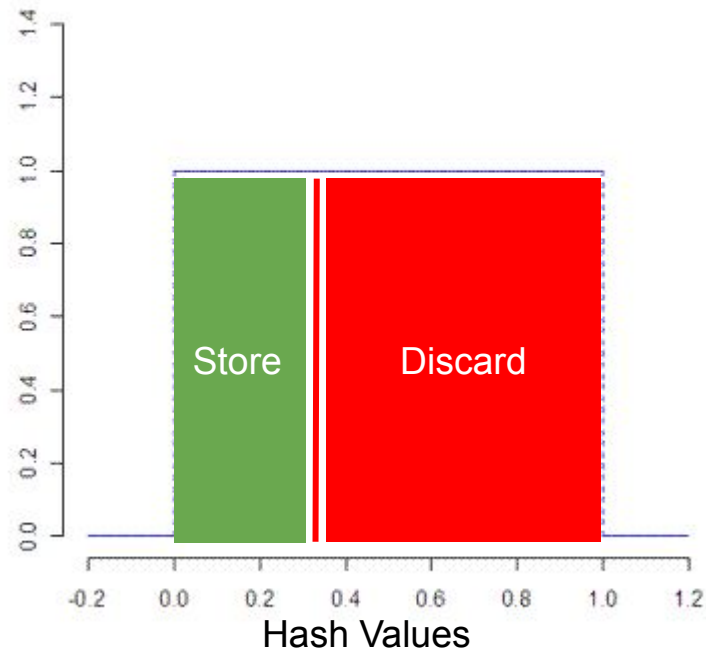
$$J(A, B) = \frac{|A \cap B|}{|A \cup B|} \approx \frac{|S(A \cup B) \cap S(A) \cap S(B)|}{|S(A \cup B)|}$$

# Calculating a Discard Threshold

- Problem:
  - Storing and discarding most k-mers
- Goal:
  - Predict which k-mers will be discarded before storage

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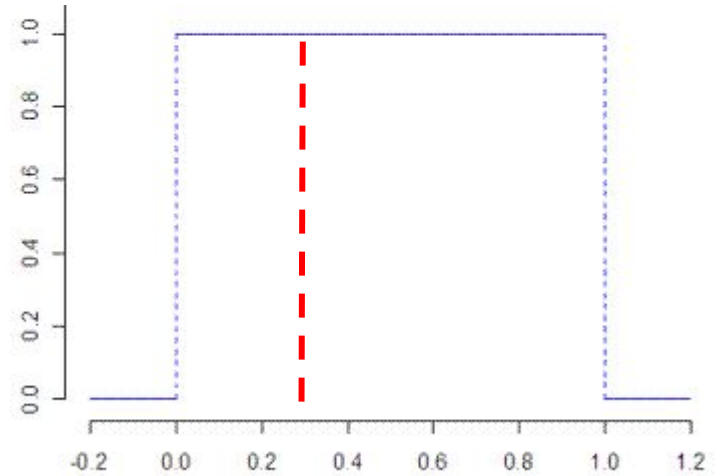




# What Threshold Value To Use?

Calculate *threshold* as *expected cutoff* for hash values

$$threshold = \frac{\text{desired \# of values}}{\text{unique values}} * \max \text{ hash value} - \min \text{ hash value} + 1$$



# Unique Values Calculation

To calculate the expected number of unique  $k$ -mers:

$$E(u) = n - n * \frac{(n - 1)^k}{n^k}$$

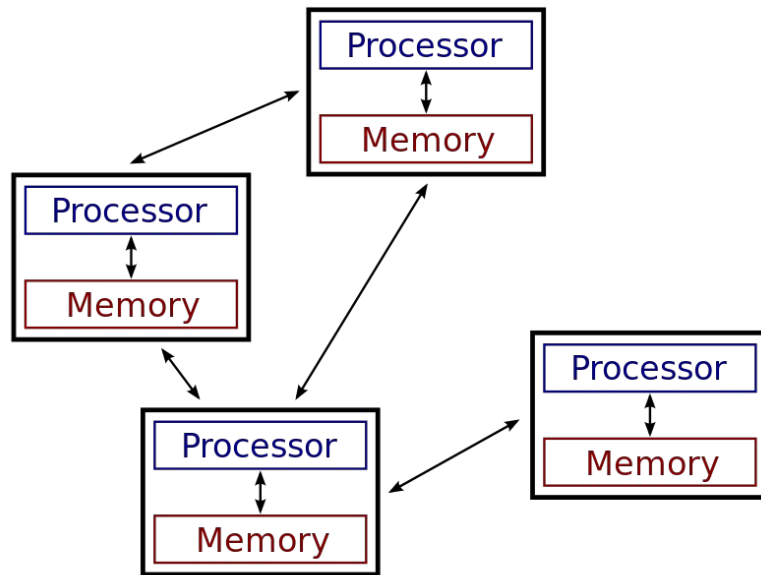
$E(u)$  = expected unique  $k$ -mers

$n$  = universe of  $k$ -mers

$k$  = number of  $k$ -mers in the file

# Distributed Computing

- Multiple processes run on separate computers
- Processes communicate via message passing (MPI)
- Communication is often run time bottleneck



# Berkeley Supercomputer



Image: <https://www.nersc.gov/news-publications/nersc-news/nersc-center-news/2016/cori-supercomputer-now-fully-installed-at-berkeley-lab/>

# Input File

File A

seq 1  
CCACACCAAAGAG...

seq 2  
GAGATTCAGCAATG...

seq 3  
CTCGAAGAGATGGA...

seq 4  
CGGCGTTAAGTTTA...

seq 5  
CGCCGATAACCCCA...

seq 6  
GACTCCGGGCTTAC...

seq 7  
ACTCTGAAAACATT...

...

seq n  
CAGCTCACCATTAC...

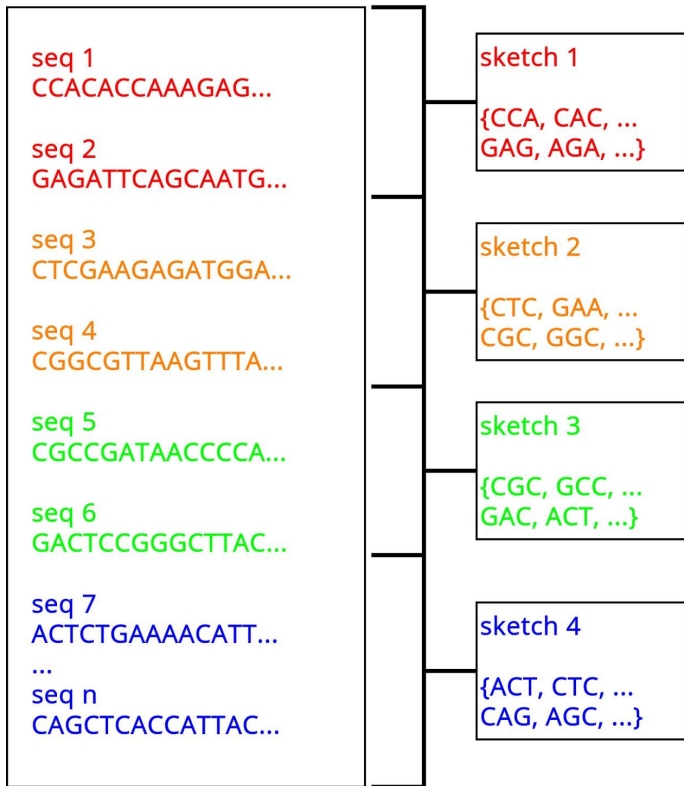


DNA Sequencing Machine

- File of reads from DNA sequencing
- May be from one or multiple organisms

# Parallel I/O using MPI

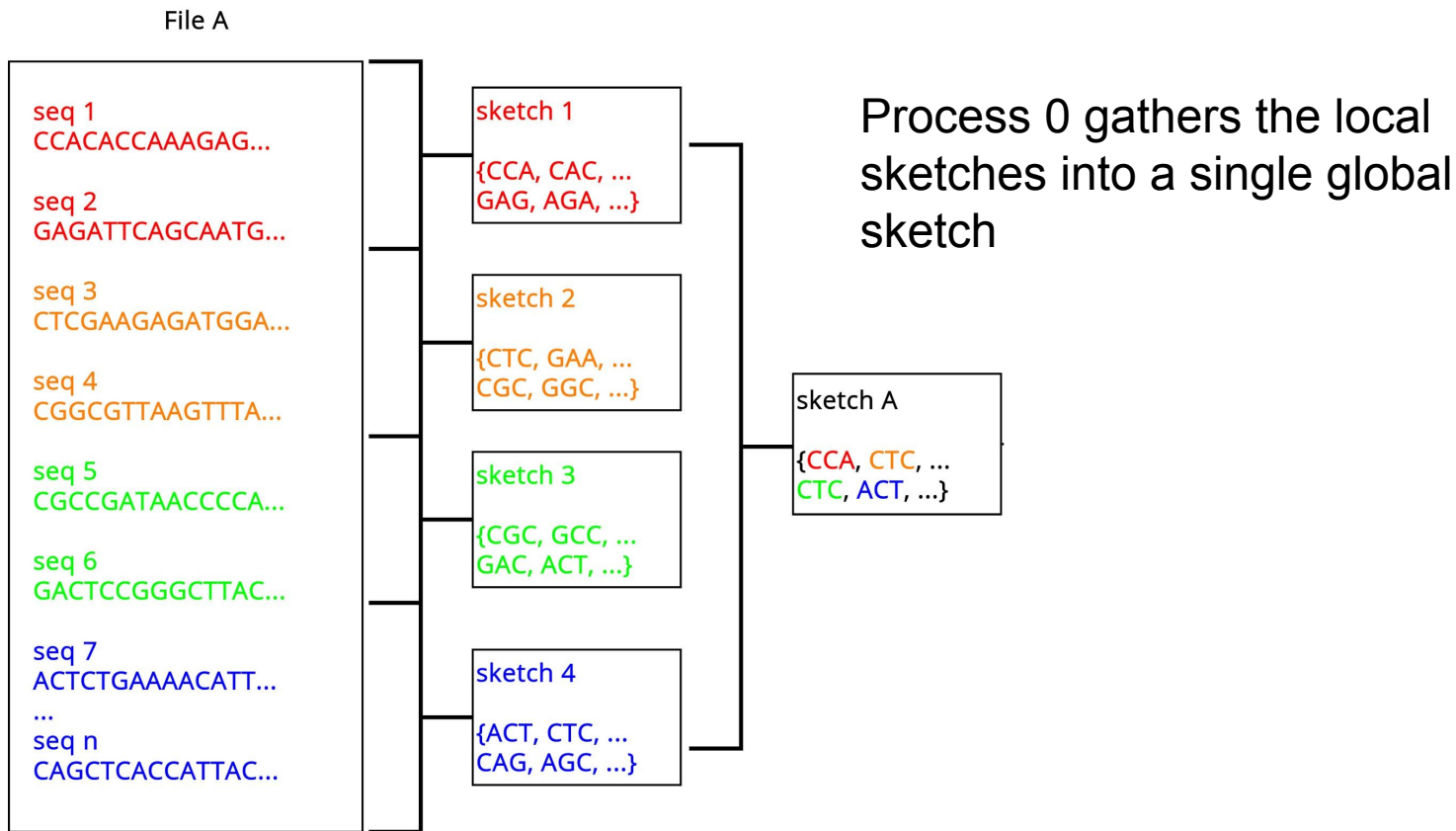
File A



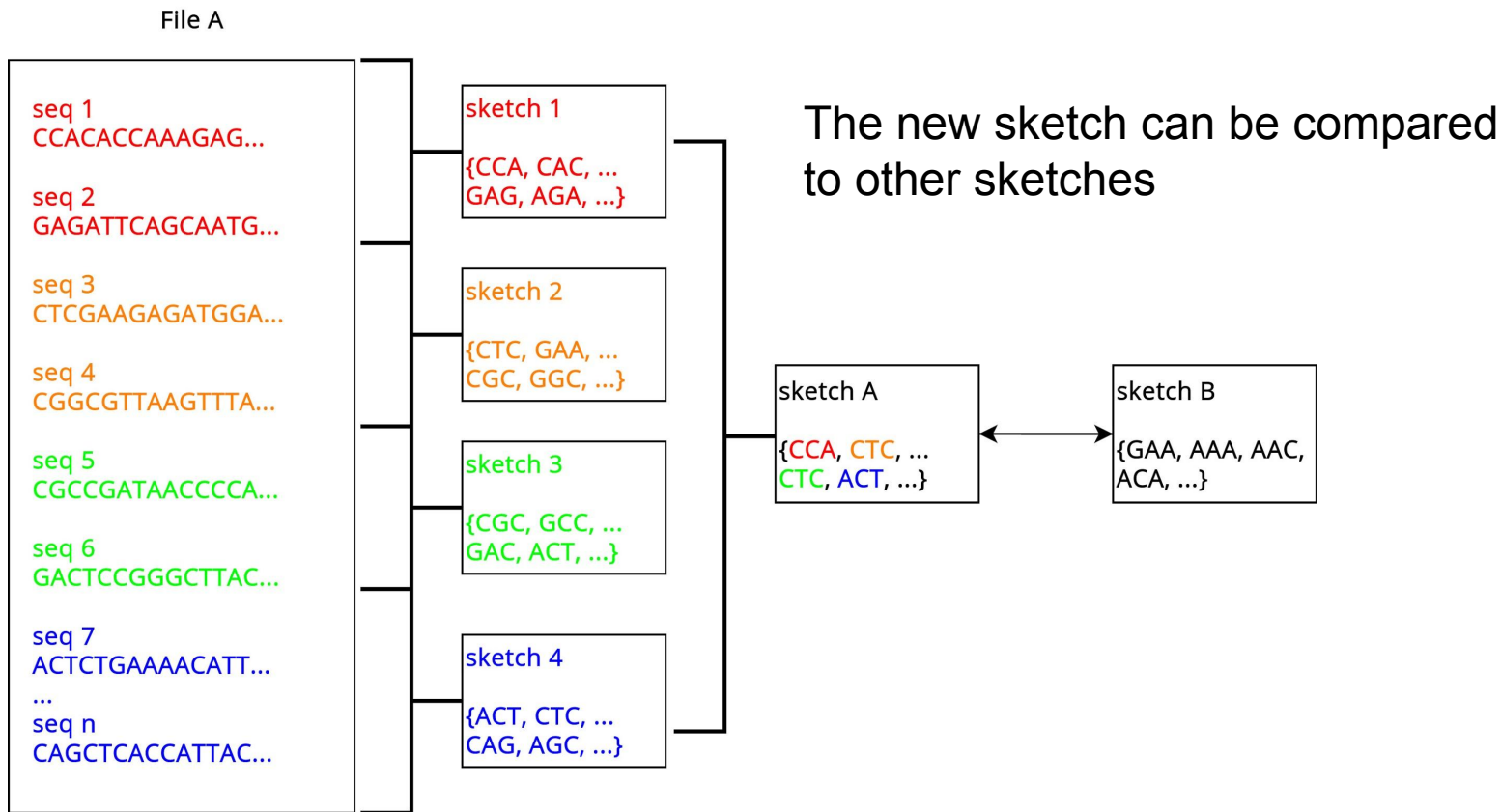
Each process simultaneously:

1. Reads a chunk of the file.
2. Creates a local sketch based on its chunk

# Aggregating Sketches



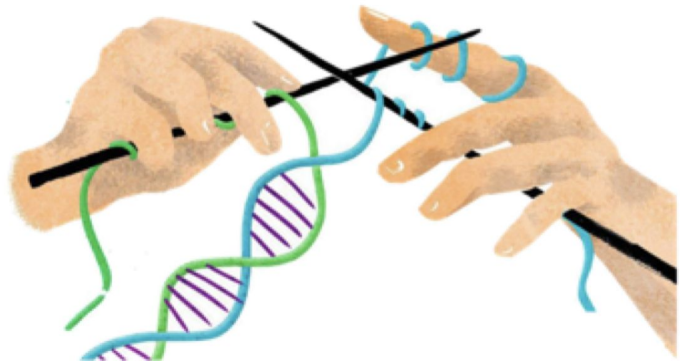
# Parallel I/O using MPI





# Conclusions

- What I've done:
  - Threshold calculation
  - Parallel I/O
- What comes next:
  - Finalizing the sketch aggregation
  - Expected efficiency for large datasets



# Thank you!

Giulia Guidi and Ben Brock  
Aydın Buluç and Kathy Yelick  
UC Berkeley

Joel Adams  
Calvin University

